advanced

Optimizing TEX code some words on speed and space

Hans Hagen PRAGMA ADE Ridderstraat 27 8061GH Hasselt NL pragma@wxs.nl

abstract

Macros can be collected in macro packages. These packages can be stored in a form that permits fast loading. Although TEX is already pretty fast, for demanding applications it makes sense to speed up TEX to the max. Switching to e-TEX and beyond is one way to achieve this, another way can be found in optimizing the macro code by means of a dedicated program. Currently the combination of both can speed up TEX runs by at least 10%.

keywords

speeding up, formats, optimizing code, e-TEX, ConTEXt

Wat is T_EX?

 T_EX is a typographic computer language with a strange character. To mention one: typographical programs and text to be typeset can be mixed. Therefore the border between programs and text is not always that clear. Take

1 2 3 \hbox spread $lem{hss4\hss}$ 5 6 7

which shows up as

 $1\ 2\ 3\ 4\ 5\ 6\ 7$

This piece of code does not really contain a program, but the next example does:

```
\setbox0=\hbox{12}\dimen0=\wd0
```

1 2 3 \hbox to $\dimen0{hss4\hss}$ 5 6 7

As long as one digit is used, the result is the same as in the previous example, because a digit has a width of .5em. It makes sense to isolate textual input and programming code, like in:

```
\def\TwoDigitsWide#1%
 {\hbox to lem{\hss#1\hss}}
```

1 2 3 \TwoDigitsWide{4} 5 6 7

The \def'd things are called macros. Quite often macros are collected into so called macro packages. That way we save ourselves the time of retyping commonly used macros and at the same time force consistency.

So, T_EX is a programming language and a program written in such a language should either be compiled into some low level machine code or interpreted at runtime. Again, T_EX is a strange breed, because it does a bit of both: T_EX the programming language is processed by T_EX the program. While read in from a file, characters and sequences of characters are translated into an internal representation and when not directly typeset, they are stored in appropriate data structures.

In our example, the definition $(\ensuremath{\mbox{Mel}}\xspace.)$ is stored for later use, and called upon while 1 2 . . is typeset.

Running T_EX

On many computer systems users invoke their prefered macro package by typing in its name at the command line and often they think that tex, latex or context is simply a program. This situation becomes even more vague when one encounters pdftex and pdflatex. (Be suspicious if you ever encounter pdfcontext, because by nature CONTEXT is not aware of any specific brand of TEX, except E-TEX!)

What actually happens is that T_EX the program is launched and instantly starts reading in the macro package one asked for. So, it's still T_EX that one's running, no matter in what way it is invoked! On many systems saying:

tex	<filename></filename>
latex	<filename></filename>
context	<filename></filename>

is just the same as:

tex &plain <filename>
tex &latex <filename>
tex &context <filename>

which is the 'official' way of calling $T_{\ensuremath{E\!} X}$ with some kind of macro package.

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Formats

The macros normally are not stored in ASCII format, but in a so called memory dump, a sort of precompiled format. One can generate such a format by saying:

```
tex --ini plain.tex \dump
tex --ini latex.ltx
tex --ini context.tex
```

When run in ini (also called virgin) mode, $T_{E}X$ reads in the macro package and when finished, it dumps its memory on disk. Such a dump normally has the suffix fmt and can be reloaded fast: <code>&plain</code> refers to the format. It does not hurt to know this! Try it, if only to be able to generate an updated format.

So, to summarize, we have T_EX the programming language, T_EX the program, and all kinds of macro packages, that can be packed into memory dumps for fast loading.

When generating an PDFT_EX format, or E-T_EX or PDF-E-T_EX format, a slightly other method applies: PDFT_EX has it own configuration information, while E-T_EX only has additional functionality when the filename is prefixed by a *. To ease the life of CONT_EXT users, there is T_EXEXEC. This PERL script wraps all brands of T_EX in a shell, so there we say:

```
texexec --make en
```

to generate a format with an english user interface. Running CONTEXT is done by saying something like:

texexec myfile
texexec --pdf myfile

So, when you're using CONTEXT, don't worry if all this formatting talk is beyond you: T_EXEXEC is there to guide you. So far for this interlude.

Tokens

 T_EX thinks in terms of tokens. Let's for this moment forget about the sometimes pretty confusing way T_EX reads and interprets the characters, and consider the next example:

Once read in memory, this example is represented in 13 tokens. The names of macros as well as the $\{, \}$ and .'s each become one internal quantity. The space after \csname ends the control sequence and does not count. I will not go into details on the amount of memory each token takes, but in general, one can say that TEX tries to minimize the amount of memory used. A macro name is stored in the so called hash table (a lookup table) and after that only references to this table are used. In the T_EX that I'm currently running, a reference to a macro name takes 8 bytes: 4 bytes for the index, a pointer to the next token, some space for an additional index (used by \charsubdef) and a rounding byte for efficient internal representation.

The main point of this short story is that once read in, T_EX does not need to translate macro names, but simply uses pointers to reach the meaning of this macro. This is not only faster, it also takes less memory. And this is why it makes sense to use memory dumps instead of reading in a macro package each time. T_EX loads faster and runs faster.

One should be aware of the fact that we're only talking of a clever way of storing data. Opposite to for instance high level programming languages like Pascal and C, no compilation is done. The meaning of the macro is stored as it is, even if its meaning is wrong in some way or another. Nearly nothing is checked and nothing is optimized. And this is one of the reasons why T_EX , while being an interpreter, is so fast.

The main reason why I could give CONTEXT a multi lingual interface without having to recode those tens of thousands lines of code, is that I used a rather high level of abstraction. Keywords and values are stored as macros. This is a sort of lucky coincidence: when CONTEXT grew, TEX's were still pretty small, so I ran out of string memory (the total length of all strings used) before I ran out of hash memory (the number of macros). So, I was forced to use macros, which after all is not that bad, because it also forced and garantees consistent use of keywords.

\def\NameKey{name}

So, name is stored once, and can be accessed by $\mbox{NameKey}$, not only many times without taking string memory, but also pretty fast. This is only true in CONTEXT deepest inners, because at the user level, one does not type in macroded keys, but verbose ones: name=hans.

- 1: \def\Name{Hans}
- 2: \def\SurName{Hagen}
- 3: \def\FullName{Hans_Hagen}

This is one way of storing names, but when one wants to save space, the next alternative is more efficient.

4: $\frac{\det[Vame]}{2}$

Counting tokens is not the way to determine this, because several types of memory are involved: the hash table, the string pool with string pointers, main memory, etc. Therefore we only really save space when more than one reference is made.

Coding keywords in macros can also be faster, especially when passing large arguments, skipping branches in conditionals and while doing certain low level lookups.

Optimizing code

One can sqeeze quite some speed from clever coding macros and using memory as efficient as possible, but when one hits the frontiers of coding itself, other methods are needed.

I started experimenting with optimization when rewriting part of the system modules. I found out that in one occasion grouping speeds up, while in another rather similar situation doing the housekeeping myself was to be prefered. Potential slow-downers are: passing long arguments, all kind of tests, especially string comparisons, list processing, and rather massive catcode changes.

Original T_EX is frozen. Named E- T_EX , its successor offers some basic programming features that are meant to speed up T_EX as well as provide more control to macro programmers. Being rather curious, I decided to change some low level code and look to what extend E- T_EX would speed up a large package like CON T_EXT . (Notice that speed was not the primary target of the E- T_EX project.)

Although currently most of the critical parts of CONT_EXT are rather well optimized —I'm still documenting, optimizing and sometimes recoding the source— I consider the measurements to be be rather representative. It is, by the way, always an interesting dilemma: do I code for speed or for readability. It's one of the reasons why font handling routines often look rather obscure: a pretty complex font mechanism demands dirty coding to get acceptable speed.

When testing for speed, it can be tempting to put some critical code in a loop, and execute this code for instance 50.000 times. Such experiments demonstrate that individual pieces of code can be rewritten in new E-TEX primitives to run much faster. Bringing down runtime from 10 seconds to 5 seconds is fine, but in practice those fragments are not executed that many times, so the gain in normal production runs is minimal. When the protection mechanisms that CoNTEXT uses for savely testing keywords are rewritten in a for E-TEX more natural way, there is even a speed penalty!

I therefore tend to conclude that it does not make much sense to recode a large macro package in $E-T_EX$ (version 2) for the sake of speed alone. This is mainly due to the fact that the critical components of CONT_EXT are already coded rather efficient. Later on I will show that in one area, $E-T_EX$ beats T_EX pretty well. Keep in mind that we are discussing speed and space. Much of $E-T_EX$'s new functionality goes beyond that and concerns better typography.

This gives us another reason to use $E-T_EX$.

An interesting observation is that using the new primitives \protected and \ifcsname saves about 500 hash entries in the current version of CONTEXT. However, making the specific pieces of macro code usefull for normal TEX and E-TEX, costs about 500 entries in normal TEX. Alas, that's the price original TEX users have to pay for progress.

Being parameter driven, CONT_EXT does a lot of string and list processing and unfortunately T_EX lacks low level support for this. When I discussed this with Taco Hoekwater, we came to the conclusion that some more straightforward support for string and list handling could speed up CONT_EXT considerably, and Taco decided to extend T_EX the program. That way we can present the E-T_EX team with well defined and tested functionality for future versions.

When testing some first versions of Taco's binaries, I wondered if it would make sense to optimize T_EX macro code in another way. To understand what I mean, I refer to a few pages back, where I introduced those hash entries and pointers. In writing macros, I try to be as clear as possible, so instead of

\cs	name	ehellc	\endcsna	me
	1	23456	7	

I code

```
\getvalue {hello}
```

and not

but, at the cost of 16 bytes overhead and a two more 'lookups':

$\frac{\text{Oifelse}}{1}$

Did you notice the difference in the number of tokens used? Using a syntax highlight editor,

\dimen0=0pt

just looks better and more readable than

\dimena\z@

Currently, and this is of course due to the fact that we are dealing with macros, we are also more talking of translating than of compiling. Some first experiments with an optimizer written in PERL were promising, but fearing unwanted side effects I decided to let this rest for a while.

The optimizer

And then Han The Thanh asked me to test his first version of PDF-E-T_EX. This merge of two rather important developments in the T_EX world: E-T_EX and PDFT_EX, drove me into making CONT_EXT more permanently E-T_EX aware as it was already pretty PDFT_EX aware. While testing, I also picked up the optimization thread.

Using the timing build into $T_EXEXEC I$ found out that on a document of average complexity, the interactive MAPS bibliography, with 700 simple pages and 50 pages of cross linked indexes and lists, rewriting some core macros to use $E-T_EX$ functionality saves about 5% run time and the optimizer gives us an additional 5%: not impressive, but useful when one considers that I quite often run jobs that take an hour or more. And, apart from more efficient coding, I expect to gain another 10% in due time.

Before I will mention some characteristics of the CONT_EXT optimizer, I need to give E-T_EX some more credit. When I first timed the test run, I found out that E-T_EX took 65% of the time the normal run needed. This proved to be due to the fact that outside E-T_EX one has to fake multiple marks, and this fake can slow down a run with many color changes on a page considerably (due to lots of list manipulations). When using normal T_EX, CONT_EXT spends half of the run time on the 50 pages mentioned before. Because marks are used for keeping track of color, and because these pages have colored hyperlinks in multiple columns, a speed penalty is paid. Unless one heavily uses color in rather complicated documents, one will probably never notice.

When color does not cross pages, which did not really happen in this document, CONTEXT can be told to switch to local color mode. In local mode, E-TEX's gain in speed is reduced to the mentioned 5%.

Some details

Back to the optimization. Apart from a few critical files, 100 modules that are part of the distribution are optimized in four converging passes (doing it in one pass is much slower). In the process, over 7.000 lines out of 80.000 lines of macro code are optimized, in many occasions, more optimizations per line. All changes are logged and some statistics are kept. Dubious and potential dangerous situations are skipped. Of course, E-TEX optimizations are optional. Don't confuse this optimization with rewriting core macros.

Some lines are skipped. Think of macro headings and calls that use delimiters like \box. Optimizing macro headings is 'not done' anyway, but the next substitutions are quite legal:

- removing redundant equal signs
- using \empty instead of {}
- using $\z \in$ instead of **0**pt
- □ changing 2, 4, 6 and 8 into constants in box primitives
- □ changing 2, 4, 6 and 8 registers into predefined ones
- using macro constants where possible
- substituting T_EX keywords by macros
- changing \getvalue cum suis into \csname
- □ optimizing all kind of \doif... macros
- \square applying E-T_EX's \ifcsname and \ifdefined
- removing redundant {} in arguments
- removing redundant T_EX keywords

As said, sometimes substituting can be dangerous. Changing for instance

\expandafter\ifx\csname...\endcsname\relax

into E-T_EX's sequence:

\unless\ifcsname...\endcsname

can lead to unwanted side effects. The pure $T_{E}X$ alternative creates a hash entry, that defaults to <code>\relax</code>, which is why we test for <code>\relax</code>. We just consider <code>\relax</code> to represent undefined. The second one does not create an entry. Consider for instance that at a certain moment, for instance in the process of font switching,

\csname...\endcsname

is expanded. This leads to ... being set to \relax. Now, when in pure T_EX , we test for existance, as expected we get reported back that the control sequence does not exist. In E- T_EX however, \ifcsname is unrelated to \relax, and therefore, E- T_EX reports that the control sequence does exist. Even if you don't completely understand what I'm talking about, you can imagine that macros that until now work perfectly ok, fail under E- T_EX .

The size of the format file (the memory dump) when optimized is currently about 75 KB less than the nonoptimized file (about 2.8 MB) and some quick tests show that another similar saving is possible. It's not that much a problem to save far more bytes, but sometimes speed goes over space:

- 1: \setbox0=\hbox{something}
- $2: \underline{\texttt{(setbox0)}}_{1 2 3 456789 11 13}$
- 3: \setbox\boxa\hbox{something} 1 2 3 456789 11 13
- 4: \setboxa\hbox{something}
- 5: \sethboxboxa{something}

The third alternative, with \boxa being \chardef'd to zero is the fastest. The last two alternatives save memory (8 bytes per macro name). It does not take much fantasy to see that the third alternative involves the two tokens \setbox and \boxa, while the fourth one takes only one (\setboxa). When resolved, this indirect reference itself takes two, therefore totalling up to three.

\def\setboxa{\setbox\boxa}

Some gain in speed in E-T_EX is due to optimizing existing T_EX code. I did not compare the results to original T_EX, but especially saving and restoring and in some situations more clever \aftergroup handling has proved to be a welcome extension.

Conclusion

What can we conclude. First that 15 years of T_EX has proven that using memory dumps makes sense. Next that for gaining some speed, developments like E- T_EX make sense too. For me however the most interesting conclusion is that some sort of preprocessing makes much sense too. Because this optimization is in many respects dependant of the way the macro package is written, it is not possible to bring this into E-T_EX the program. Anyhow, given this 10% speed gain, makes me very optimistic about Taco's estimations on string and list processing. It can also be a step towards a higher level of typographic programming in T_EX, where more readable definitions are compiled into their lower level T_EX equivalents.

Literature

- □ The E-T_EX manual, version 2, February 1998, Peter Breitenlohner.
- E-T_EX, a 100% compatible successor to T_EX, Philip Taylor, EuroT_EX 1995 proceedings.
- Examples: especially the CONTEXT modules syst-gen, supp-mrk, font-ini.
- Unpublished e-mails, Taco Hoekwater.